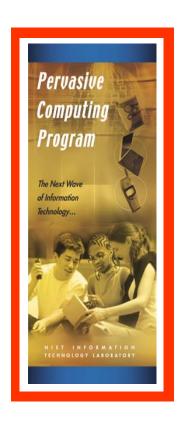
## Understanding and Improving Service Discovery in Hostile and Volatile Environments



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NRC Review Panel February 8, 2002

## Dynamic Discovery Protocols in Essence

Dynamic discovery protocols enable *network elements* (including software clients and services, as well as devices):

- (1) to *discover* each other without prior arrangement,
- (2) to express opportunities for collaboration,
- (3) to *compose* themselves into larger collections that cooperate to meet an application need, and
- (4) to detect and adapt to changes in network topology.

## Selected First-Generation Dynamic Discovery Protocols



#### Presentation Outline

- Project objective and plan (one slide)
- Accomplishments over past year (one slide)
- Context setting (two slides)
  - Alternative architectures for service-discovery systems
  - Technical approach to modeling and analysis in first phase of the project
- Understanding consistency-maintenance in service-discovery systems under increasing communication failure (12 slides)
- Summary of progress to date (one slide)
- Plans for upcoming year (one slide)
- Addenda
  - Additional details on consistency-maintenance experiments
  - Issues relayed to Sun after correctness analysis of Jini specification
  - Performance measurement results from Jini and UPnP implementations

### **Project Objective**

Research, design, and characterize self-adaptive mechanisms to improve performance of discovery protocols in hostile and volatile environments.

### Project Plan in Three Phases

<u>Phase I</u> – characterize performance of selected service-discovery protocols (Universal Plug-and-Play – UPnP – and Jini) as specified and implemented

- develop architectural models for each protocol
- establish performance benchmarks based on default or recommended parameter values and on required or most likely implementation of behaviors

<u>Phase II</u> – design, simulate, and characterize self-adaptive mechanisms to improve performance in volatile environments

- devise algorithms to adjust control parameters and behavior in each protocol
- simulate performance of each algorithm against benchmark performance
- select most promising algorithms for further development

<u>Phase III</u> – implement and validate the most promising algorithms in publicly available reference software (from Sun Microsystems and Intel)

### Accomplishments This Year

#### **Modeling**

- Completed Rapide architectural-description language model of UPnP specification (Jini specification completed last year)
- Developed framework of (SLX) discrete-event simulation model for UPnP (model patterned after publicly available Intel Linux implementation)

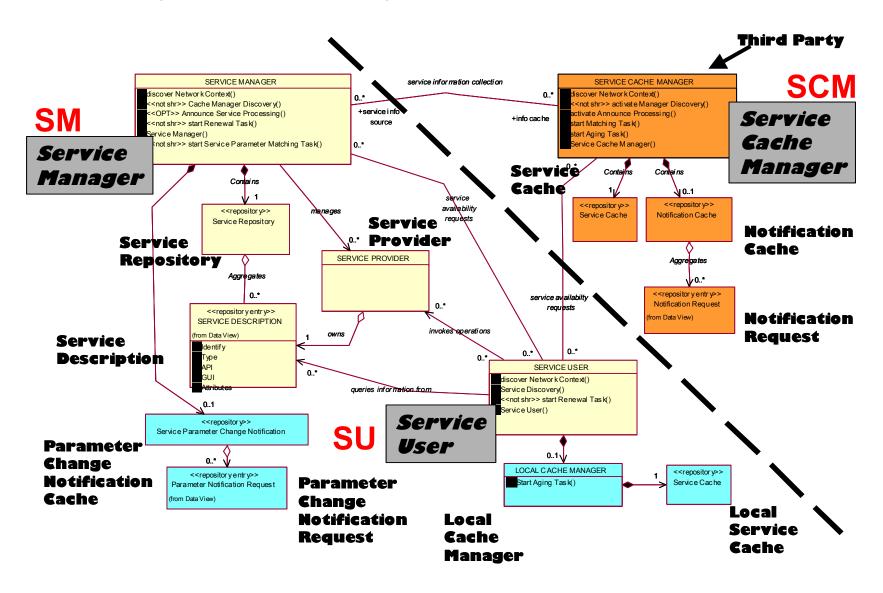
#### **Analysis**

- Explored correctness of Jini model against various consistency conditions; found four areas of concern; briefed Sun's Jini design team on findings (see addendum for more information)
- Investigated consistency maintenance in Jini and UPnP models under increasing interface-failure rate (technical subject of current talk)

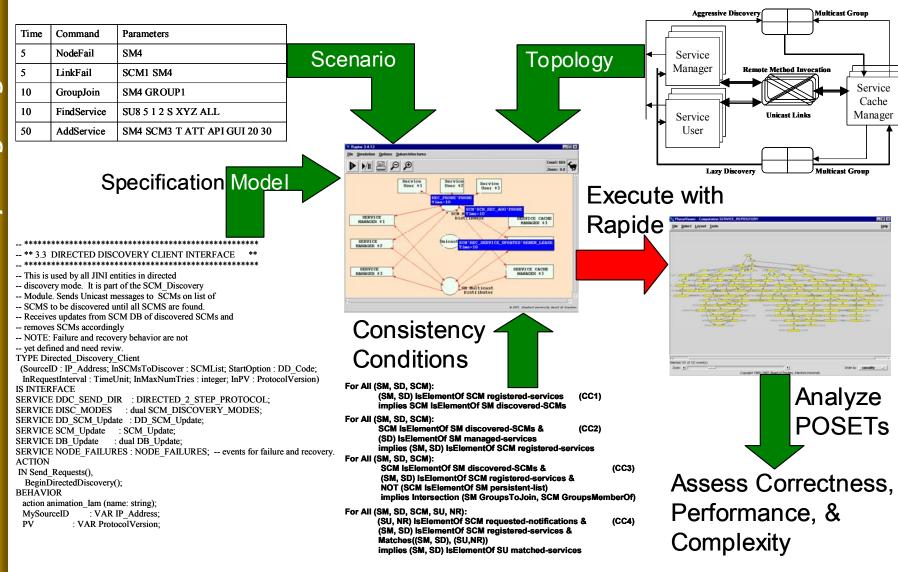
#### Measurement (see addendum for more information)

- Designed independent benchmark service and scenarios for servicediscovery systems
- Created synthetic workload generation tools for emulating the behavior of large scale, dynamic, ad-hoc networking environments
- Conducted baseline performance measurement experiments to compare the performance of UPnP and Jini in benchmark scenarios

## Two Party vs. Three Party Architectures



## Phase I Modeling Approach: Use Rapide to Model and Understand Dynamics of Jini and UPnP



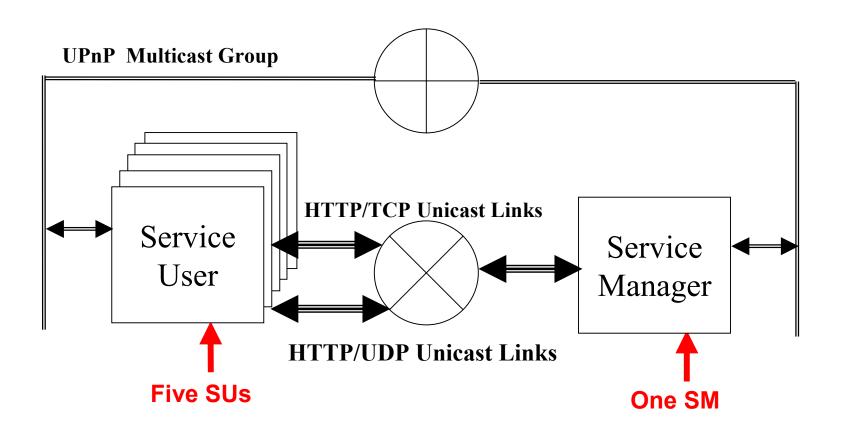
## Understanding Consistency Maintenance in Discovery Architectures during Interface Failure

How do various service discovery architectures, topologies, and fault-recovery mechanisms perform under deadline during interface failure?

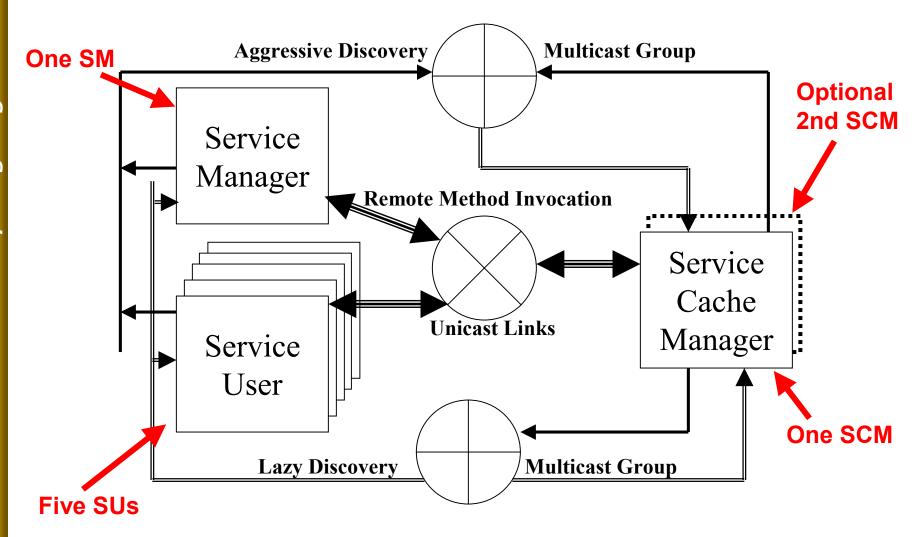
### **Outline of Experiment**

- Deploy models of two-party (UPnP) and three-party (Jini) architectures with one SM and five SUs (for Jini include two topologies – one and two SCMs).
- Ensure initial discovery and information propagation completed.
- Introduce a change in the service description at the SM, and establish a deadline for propagating the new information to all SUs.
- Measure the number of messages exchanged and the latency required to propagate the information to all SUs, or until the deadline arrives, under two different propagation mechanisms: polling and eventing.
- Repeat this experiment while varying the percentage of interface failure time for each node up to 75% (in increments of 5%).

## Two-Party (UPnP) Topology for Experiment



### Three-Party (Jini) Topologies for Experiment



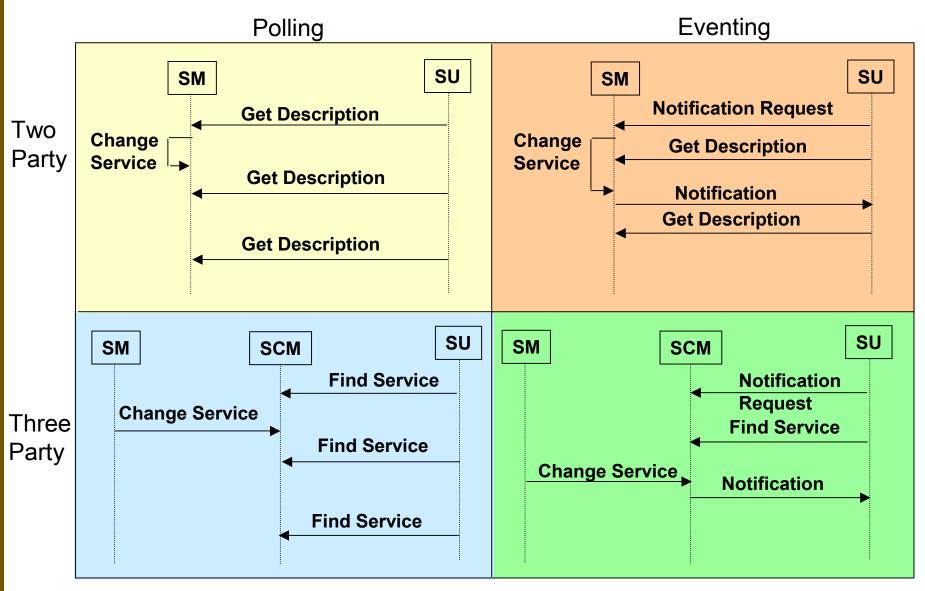
### Failures Interfere with Discovery and Information Propagation

- Interface Failure Tx, Rx, or Both
  - Due to nearby enemy jamming or other interference
  - Due to multi-path fading during mobility
- Path Loss Pt-Pt or Area-Area and Full-duplex or Half-duplex
  - Due to persistent congestion
  - Due to physical link cuts
  - Due to enemy jamming at routers
- Message Loss under both UDP and TCP (>delay)
  - Due to sporadic or distant enemy jamming or other interference
  - Due to transient congestion
  - Due to multi-path fading during mobility
- Node Failure Partial or Complete with variable persistence of information
  - Due to enemy bombardment or cyber attacks
  - Due to mobility associated with military operations

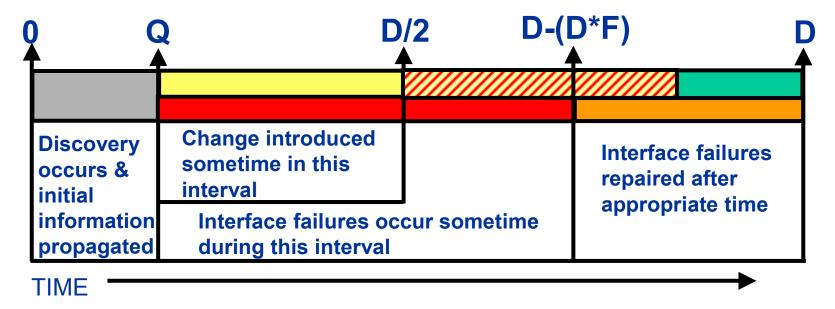
## Discovery Systems Divide Recovery Responsibilities: Lower Layers, Discovery Protocol, and Application

- Selective Reliable Delivery by Lower Layers
  - TCP attempts retransmissions (basis for Jini-RMI and UPnP-HTTP/TCP)
  - UDP messages in UPnP sent as multiple copies
- Periodic Announcements by Discovery Protocol
  - Allows caching nodes to discard information when TTL expires
  - Jini includes aggressive search at node start up, while UPnP permits nodes to undertake aggressive search at any time
- Periodic Refreshing of Resources Required by Discovery Protocol
  - Allows resource owner to free resource when refresh period expires
- Remote Exceptions Issued by Protocol Over TCP Links
  - Allows application to take recovery action: Ignore? Retry? Discard knowledge of service or resource?

## Consistency-Maintenance Mechanisms for Experiment



## Interface-Failure Model for Experiment



- 1. Choose a time to introduce the change [uniform(Q, D/2)]
- Random Processes
- 2. For each node, choose a time to introduce an interface failure [uniform(Q, D-(D\*F))]
- 3. When each interface failure occurs, choose the scope of the failure, where each of [Rx, Tx, Both] has an equal probability

Q = end of quiescent period (100 s in our experiment)

D = propagation deadline (5400 s in our experiment)

F = Interface Failure Rate (variable from 0% - 75% in 5% increments in our experiment)

### Our Modeled Responses to Remote Exceptions

("Approximate" – see addendum for details)

- Ignore REX Received
  - When replying to a remote-method invocation
  - When attempting to cancel a lease
  - When attempting to renew a lease (But then attempt to obtain a new lease)
- Retry Operation for Some Period of Time Then Quit If Not Successful
  - When attempting to register for notification events
  - When a UPnP SU requests service descriptions
  - When Jini SM attempts to register service or change service on SCM
- Eliminate Local Knowledge of Discovered Entity
  - When remote exceptions occur for too long a period
  - When UPnP SU fails to hear scheduled SM announcement

### Metrics Devised for Consistency-Maintenance Experiments

(see addendum for formal definitions)

#### Update Responsiveness (R)

Assuming that information is created at a particular time and must be propagated by a deadline, then the difference between the deadline and the creation time represents available time in which to propagate the information. Update Responsiveness, *R*, measures the *proportion of the available time remaining after the information is propagated*. [1 = all time remains and 0 = no time remains]

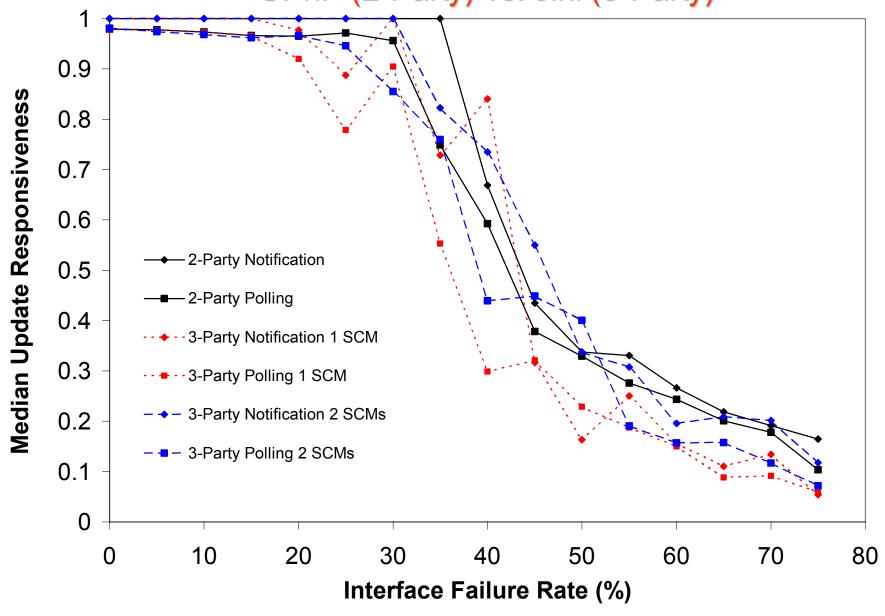
#### • Update Effectiveness (*U*)

Update Effectiveness, *U*, measures the *probability that information will propagate* successfully to a SU before some deadline, D. [1 = information will be propagated and 0 = information will not be propagated]

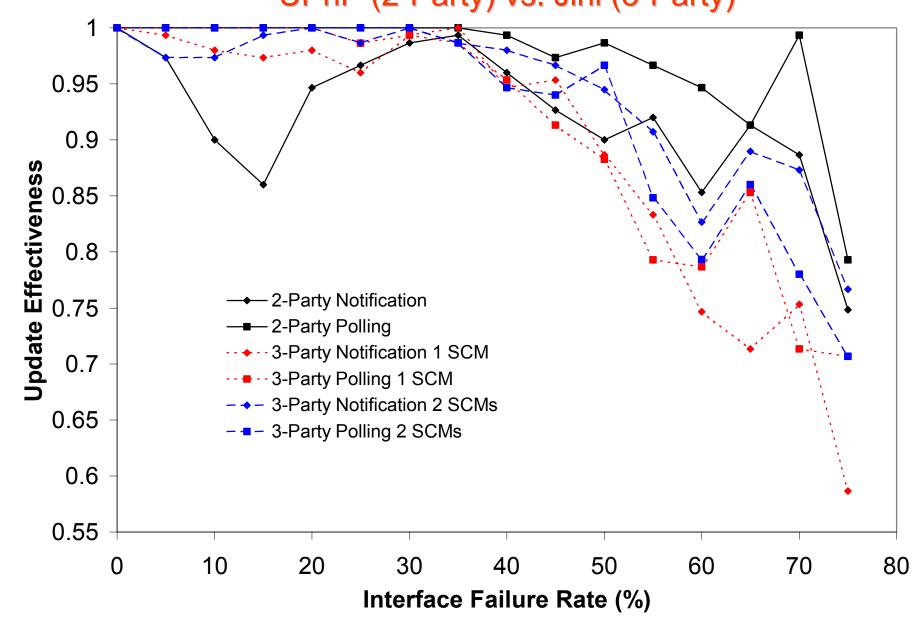
#### Update Efficiency (E)

Given a specific topology of SUs and SMs in a discovery system, examination of the available architectures (two-party and three-party) and mechanisms (polling and eventing) will reveal a minimum number of messages that need to be sent to propagate information from all SMs to all SUs in the topology. Update Efficiency, *E*, can be measured as the *ratio of the minimum number of messages needed to the actual number of messages observed*. [1 = only minimum number of messages needed]

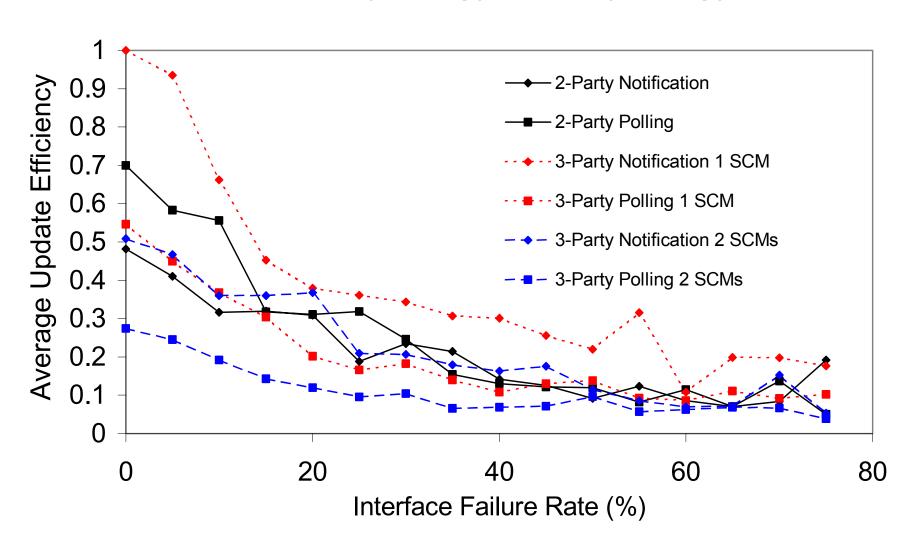
# Median Update Responsiveness UPnP (2-Party) vs. Jini (3-Party)



# Update Effectiveness UPnP (2-Party) vs. Jini (3-Party)



## Average Update Efficiency UPnP (2-Party) vs. Jini (3-Party)



### Summary of Progress to Date

- Models of service-discovery protocols (written in the Rapide architecturedescription language):
  - Jini (models the Sun Microsystems version 1.1 specification)
  - UPnP (models the Microsoft version 1.0 specification)

#### Papers

- "Analyzing Properties and Behavior of Service Discovery Protocols Using an Architecture-based Approach", C. Dabrowski and K. Mills, Proceedings of Working Conference on Complex and Dynamic Systems Architectures
- "Understanding Consistency Maintenance in Service Discovery Architectures during Communication Failure", C. Dabrowski, K. Mills, and J. Elder, submitted to the 3<sup>rd</sup> International Workshop on Software Performance.
- "Survivable Software for Harsh Environments", C. Dabrowski, K. Mills, and J. Elder, submitted to MILCOM 2002.

#### Interactions

- Gave an invited "Tech-a-Tech" presentation to Sun Jini team
- Presented work at 2001 Pervasive Computing Conference
- Obtained co-funding from DARPA and ARDA

### Plans for Upcoming Year

#### **Modeling**

- Complete (SLX) discrete-event simulation model for UPnP
- Validate model against current results from Rapide model
- Develop SLX simulation models for Jini and SLP
- Extend generic structural model of service discovery protocols to include message exchanges and consistency conditions, and publish the extended generic model

#### **Analysis**

- Investigate consistency maintenance in Jini and UPnP models under increasing message-loss rate
- Investigate consistency maintenance in Jini and UPnP models under node and thread failure
- Propose, model, and evaluate selected self-adaptive mechanisms for UPnP
- Consider analytical model for update effectiveness during interface failure

## Slides Containing Additional Details on Consistency-Maintenance Experiments

## Equating a Generic Structural Model of Service Discovery Architectures to Selected Commercial Discovery Systems

Generic Model	Jini	UPnP	SLP
Service User	Client	Control Point	User Agent
Service Manager	Service or Device Proxy	Root Device	Service Agent
Service Provider	Service	Device or Service	Service
Service Description	Service Item	Device/Service Description	Service Registration
Identity	Service ID	Universal Unique ID	Service URL
Туре	Service Type	Device/Service Type	Service Type
Attributes	Attribute Set	Device/Service Schema	Service Attributes
User Interface	Service Applet	Presentation URL	Template URL
Program Interface	Service Proxy	Control/Event URL	Template URL
Service Cache Manger	Lookup Service	not applicable	Directory Service Agent (optional)

## The Six Combinations of Architecture, Topology, and Consistency-Maintenance Mechanism Used in Experiments

Architectural Variant	Protocol Basis	Consistency-Maintenance Mechanism
Two-Party	UPnP	Polling
Two-Party	UPnP	Notification (with notification registration on SM)
Three-Party (Single SCM)	Jini	Polling (with service registration on SCM)
Three-Party (Single SCM)	Jini	Notification (with service registration and notification registration on SCM)
Three-Party (Dual SCM)	Jini	Polling (with service registration on SCM)
Three-Party (Dual SCM)	Jini	Notification (with service registration and notification registration on SCM)

## Specific Division of Failure-Recovery Responsibilities Used in Experiments

Responsible Party	Recovery Mechanism	Two-Party Architecture (UPnP)	Three-Party Architecture (Jini)		
Lower-Layer UDP		No recovery	No recovery		
Protocols T(	TCP	Issue REX in 30-75 s	Issue REX in 30-75 s		
Lazy Discovery		SM: announces with <i>n</i> (3+2 <i>d</i> + <i>k</i> ) messages every 1800 s	SCM: announces every 120 s		
Discovery Protocols	A	SU: issues <i>Msearch</i> every 120 s	SU and SM: issue seven probes		
1 100000	Aggressive Discovery	(after purging SD)	(at 5 s intervals) only during startup		
	Ignore REX	SU: HTTP Get Poll	SU: FindService Poll		
Application Software		SM: Notification	SCM: Notification		
	Retry after REX	SU: HTTP Get after discovery retry in 180 s (retries ≤ 3)  Subscribe requests retry in 120s	SM: depositing or refreshing SD copy on SCM retry in 120s SU: registering and refreshing notification requests with SCM retry in 120 s		
	Discard Knowledge	SU: purge SD after failure to receive SM announcement within 1800 s	SU and SM: purge SCM after 540 s of continuous REX		

## Significant Parameters and Values Used in Experiments

'	Parameter	Value		
	Polling interval	180 s		
Behavior in both two- and three-	Registration TTL	1800 s		
party architectures	Polling interval Registration TTL Time to retry after REX (if applicable) Announce interval Msearch query interval SU purges SD Probe interval Announce interval SM or SU purges SCM Afte Failure incidence Once Failure scope Tra both Failure duration 5%  UDP transmission delay TCP transmission delay	120 s		
UPnP-specific	Announce interval	1800 s		
behavior for two-	Msearch query interval	120 s		
party architecture	SU purges SD	At TTL expiration		
Jini-specific	Probe interval	5 s (7 times)		
behavior for three-	Announce interval	120 s		
party architecture	SM or SU purges SCM	After 540 s with only REX		
	Failure incidence	Once per run for each node		
Jini-specific behavior for three- party architecture Interface failure parameters	Failure scope	Transmitter, receiver, or both with equal likelihood		
paramotoro	Failure duration	5% increments of 5400 s from 0 to 75%		
		10 us constant		
Transmission and	TCP transmission delay	10-100 us uniform		
processing delays	Per-item processing delay	100 us for cache items 10 us for other items		

### Console Output from a Sample Experiment Run

```
Rate - 5
Run number - 21
SM 1 OUT Interface
                       down 365, up 635
SCM 1 OUT Interface down 2417, up 2687
SCM 2 IN & OUT Interface down 519, up 789
SU 1 IN Interface
                       down 2238, up 2508
                       down 3256, up 3526
SU 2 IN Interface
SU 3 IN Interface
                       down 207, up 477
                      down 2876, up 3146
SU 4 OUT Interface
SU 5 IN Interface
                      down 4478, up 4748
```

#### Performance:

```
SM 1 346.00000 346.00000 6 17

SCM 1 346.00000 346.00016 61 102

SCM 2 346.00000 346.00015 61 105

SU 1 346.00000 346.00109 0 11

SU 2 346.00000 346.00109 0 11

SU 3 346.00000 5400.00000 4 11

SU 4 346.00000 346.00109 0 11

SU 5 346.00000 346.00114 0 11
```

## Update Responsiveness (R)

Let *D* be a deadline by which we wish to propagate information to each service user (SU) node (*n*) in a service discovery topology.

Let  $t_C$  be the creation time of the information that we wish to propagate, where  $t_C < D$ .

Let  $t_{U(n)}$  be the time that the information is propagated to SU n, where n = 1 to N, and N is the total number of SUs in a service discovery topology.

Define information-propagation latency (L) for an SU n as:

$$L_n = (t_{U(n)} - t_C)/(\max(D, t_{U(n)}) - t_C).$$

Define update responsiveness (R) for an SU n as:

$$R_n = 1 - L_n$$
.

### Update Effectiveness (U)

Let the definitions related to Update Responsiveness, R, hold.

Let X represent the number of runs during which a particular service discovery topology is observed under identical conditions.

Recalling that *N* is the total number of SUs in a service discovery topology, define the number of SUs observed under identical conditions as:

$$O = X \cdot N$$
.

Define the probability of failure to propagate information to an SU as:

$$P(F) = (\text{count}(R_{i,j} == 0))/O$$
, where  $i = 1..N$  and  $j = 1..X$ .

Define the Update Effectiveness for a given set of conditions as:

$$U=1-P(F).$$

## Update Efficiency (E)

Let the preceding definitions associated with Update Responsiveness and Update Effectiveness hold.

Let *M* be the minimum number of messages needed to propagate information from all SMs to all SUs.

Let S be the observed number of messages sent while attempting (failures may occur) to propagate information from all SMs to all SUs in a given run of the topology.

Define average Update Efficiency as:

$$E_{avg} = (\text{sum}(M/S_k))/X$$
, where  $k = 1..X$ .

## Summary Statistics for Performance of Each Combination on Each Metric

•	Mean (across all interface-failure rates)				
	Median Responsiveness	Effectiveness	Average Efficiency		
Two-Party Notification	0.663	0.921	0.212		
Two-Party Polling	0.615	0.973	0.251		
Three-Party Notification (Single SCM)	0.601	0.894	0.389		
Three-Party Polling (Single SCM)	0.530	0.911	0.201		
Three-Party Notification (Dual SCM)	0.655	0.942	0.221		
Three-Party Polling (Dual SCM)	0.587	0.927	0.110		

## 95% C.I. for Each Metric-Combination at Selected Failure Rates

•	Responsiveness		Ef	Effectiveness		Efficiency			
_	5%	40%	75%	5%	40%.	75%	5%	40%	75%
Two-Party Notification	1.000	0.561	0.111	0.970	0.954	0.709	0.354	0.065	0.031
	1.000	0.783	0.162	0.977	0.966	0.787	0.467	0.220	0.354
Two-Party Polling	0.975	0.501	0.076	1.000	0.993	0.760	0.501	0.031	0.042
	0.980	0.849	0.138	1.000	0.993	0.826	0.666	0.230	0.059
Three-Party Notification (Single SCM)	1.000	0.605	0.042	0.993	0.939	0.521	0.827	0.099	0.033
	1.000	1.000	0.095	0.993	0.955	0.652	1.000	0.504	0.320
Three-Party Polling (Single SCM)	0.974	0.244	0.043	1.000	0.946	0.660	0.387	0.043	0.040
	0.980	0.412	0.083	1.000	0.960	0.753	0.512	0.173	0.164
Three-Party Notification (Dual SCM)	1.000	0.562	0.099	0.970	0.977	0.730	0.335	0.035	0.009
	1.000	1.000	0.143	0.977	0.983	0.803	0.599	0.290	0.096
Three-Party Polling	0.974	0.391	0.056	1.000	0.939	0.660	0.218	0.033	0.019
(Dual SCM)	0.986	0.543	0.096	1.000	0.955	0.753	0.273	0.103	0.059

## Slides Outlining Four Issues Shared With Sun Based on Exploring Correctness of our Jini Model

### SUs Registered for Notifications Should Receive Them

#### Sample Consistency Condition #4 (Service Notification)

For All (SM, SD, SCM, SU, NR):

(SU, NR) IsElementOf SCM requested-notifications & (SM, SD) IsElementOf SCM registered-services & Matches ((SM, SD), (SU, NR))

implies (SM, SD) IsElementOf (SU matched-services)\*

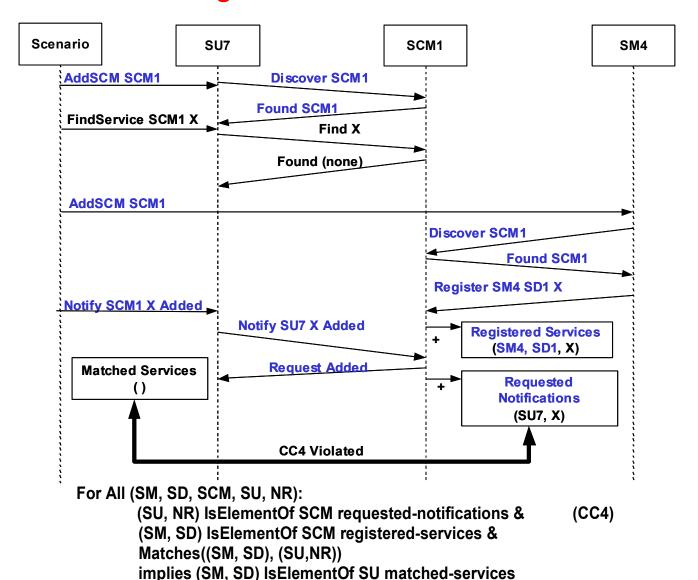
...that is, if an SU has requested notification with a Service Cache Manager of a service that matches a service description registered by a Service Manager on the same Cache Manager, then that service description should be provided to the Service User.

\*Assuming absence of network failure and normal delays due to updates

- SM is Service Manager
- SD is Service Description
- SCM is Service Cache Manager
- SU is Service User
- NR is Notification Request

- requested-notifications is a set of (SU,NR) pairs maintained by the SCM
- registered-services is a set of (SM,SD) pairs maintained by the SCM
- matched-services is the set of (SM,SD) pairs maintained by the SU

## Should the Jini Specification Advise about Possibility for Registration Race Condition?



## Discovered SCMs Should Be Members of Same Group As Discovering Entity

Sample Consistency Condition #3 (Intersecting Group Membership)

For All (SM, SD, SCM):

SCM IsElementOf SM discovered-SCMs
(SM, SD) IsElementOf SCM registered-services
NOT (SCM IsElementOf SM persistent-list)
implies Intersection (SM GroupsToJoin, SCM GroupsMemberOf)\*

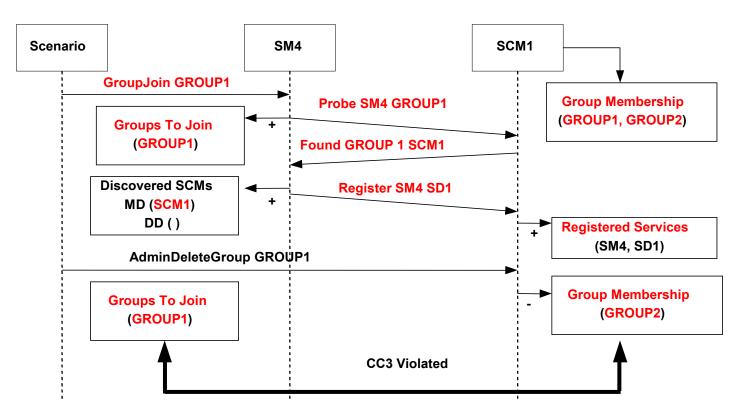
...that is, if a Service Manager has discovered, and registered its service descriptions on, a Service Cache Manager that is not on the Service Manager's persistent list, then the Service Manager must be seeking group membership in at least one group the Service Cache Manager belongs to.

\*Assuming absence of network failure and normal delays due to updates

- SM is Service Manager
- SD is Service Description
- SCM is Service Cache Manager

- registered-services is a set of (SM,SD) pairs maintained by the SCM
- discovered-SCMs is a set of SCMs discovered by the SM
- Persistent-list is the set of SCMs the SM is seeking though directed discovery

# What Might Happen When SCM Changes Group Membership Dynamically?



For All (SM, SD, SCM):

SCM IsElementOf SM discovered-SCMs & (CC3) (SM, SD) IsElementOf SCM registered-services & NOT (SCM IsElementOf SM persistent-list) implies Intersection (SM GroupsToJoin, SCM GroupsMemberOf)

### Jini Entities Should Only Register with Discovered SCMs

### Sample Consistency Condition #1(Register Only with Discovered SCMs)

For All (SM, SD, SCM): (SM, SD) IsElementOf SCM registered-services implies SCM IsElementOf SM discovered-SCMs\*

...that is, a Service Manager should register its Services on an Service Cache Manager <u>only if</u> it maintains that Cache Manager on its "known SCM" LIst.

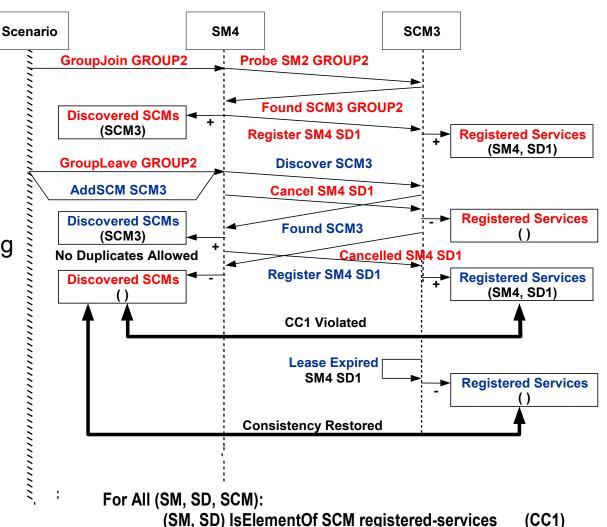
\*Assuming absence of network failure and normal delays due to updates

- SM is Service Manager
- SD is Service Description
- SCM is Service Cache Manager

- registered-services is a set of (SM,SD) pairs maintained by the SCM
- discovered-SCMs is a set of SCMs discovered by the SM

## Could the Jini Specification Lead to Implementations Exhibiting Undesired Interaction between Directed and Multicast Discovery?

Based on one possible interpretation of specification using a single-list assumption



implies SCM IsElementOf SM discovered-SCMs

# Jini Entities Should Register with All Discovered SCMs about which they maintain knowledge

Sample Consistency Condition #2 (Register Service with Discovered SCMs)

For AII (SM, SD, SCM):

SCM IsElementOf SM discovered-SCMs & SD IsElementOf SM managed-services implies (SM, SD) IsElementOf SCM registered-services\*

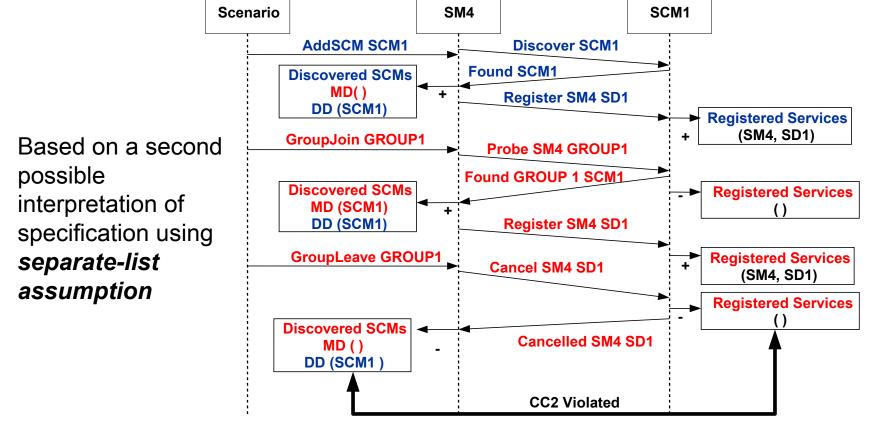
...that is, a Service Manager should register its Services on an Service Cache Manager <u>if</u> the Service Manager has discovered the Cache Manager and is maintaining the SCM identifier on its "known SCM" LIst.

\*Assuming absence of network failure and normal delays due to updates

- SM is Service Manager
- SD is Service Description
- SCM is Service Cache Manager

- discovered-SCMs is a set of SCMs discovered by the SM
- managed-services is a set of (SM,SD) pairs maintained by the SM
- registered-services is the set of (SM, SD) pairs maintained by the SCM

## Could the Jini Specification Lead to Implementations Exhibiting Undesired Interaction between Directed and Multicast Discovery?



For All (SM, SD, SCM):

SCM IsElementOf SM discovered-SCMs & (CC2)

(SD) IsElementOf SM managed-services

implies (SM, SD) IsElementOf SCM registered-services

# Slides Describing Performance Measurement Results from Jini and UPnP Implementations

### Measurement Methodology

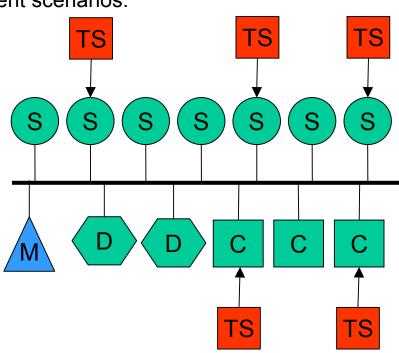
- SDP Architectures & Protocols Differ Considerably
  - Directory based vs flat peer-to-peer (combinations)
  - Java RMI and Serialized Objects vs. HTTP/SOAP/GENA and XML
- How to reasonably compare the performance of such diverse technologies?
- Usage-Based Scenarios & Metrics
  - Service initiation restart, auto configuration, advertisement, renewal
  - Client active query restart, by name, single instance, multiple instances, all instances
  - Client passive monitoring persistent query for new instances, network restart
  - Event Notification and control registration latency, notification latency, distributed control performance (control + event notification).
- SDP-Independent Benchmark Service
  - Simple counting device/service that can be used to exercise all significant discovery/control capabilities of Jini and UPnP.
  - Supports get/set service, GUIDs/attributes/type, control, event notification, GUI
- Implementation-Independent Measurement Tools.
  - Measure on-the-wire
  - Response/Load

### Synthetic Workload Generation Tools

- Objective Emulate large, dynamic environments of 100's of devices/services and 10's of control points / clients.
  - Dynamic devices providing the benchmark service.
  - Scripted control points execute measurement scenarios.

### SDP Experimenters Toolkits

- -Drive real SDP implementations
- Emulate the behavior of a large number of dynamic devices
- –Emulate the behavior control points/ scripted behavior for testing
- Jini & UPnP Initial development complete
  - SunMS Jini, Intel UPnP on Linux platforms.
  - Target of 100's of devices and 10's of control points met.



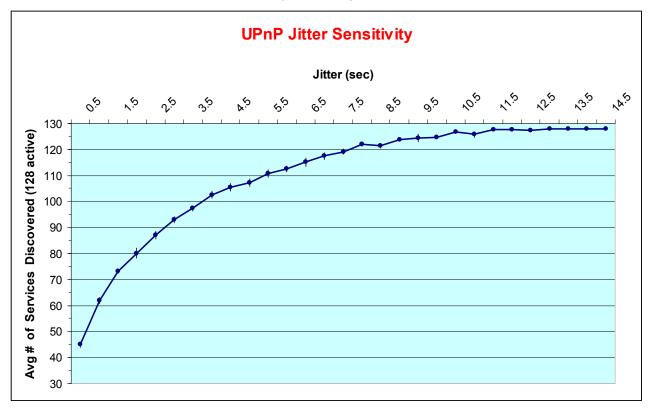
### Sample Measurement Results

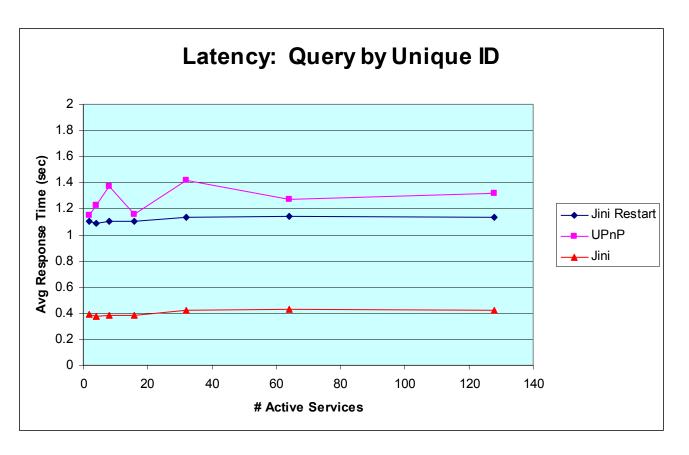
#### **Notes:**

- Example results of Client Active Query scenarios.
- Query experiments measure latency / load through the client download of the service proxy / description.
- Results show two scenarios for Jini
  - Jini Restart includes overhead of client discovery and first access of look up server.
  - Jini query from "warm" client, after lookup server Discovery has completed.
- Default UPnP jitter value is 2.5 seconds
  - In device poll experiment, jitter values of 5.5 sec (64 services) and 11 sec (128 services) were used.

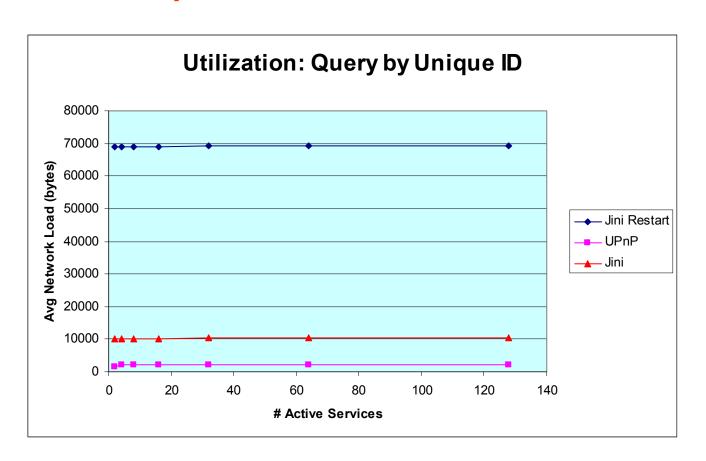
### **Linux UPnP Scaling Problems**

- Problems encountered in achieving initial scaling goals for device emulation tools.
  - UPnP scalability above 40 devices a function of protocol tuning parameters (e.g., response jitter, multicast retransmission factor).
  - Errors in implementation of jitter algorithms

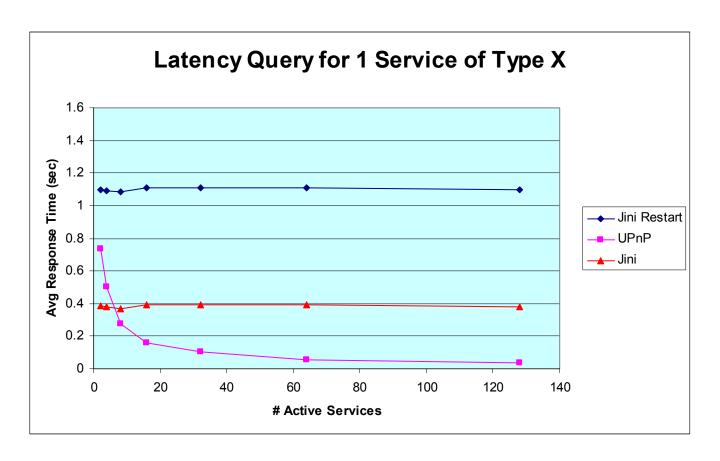




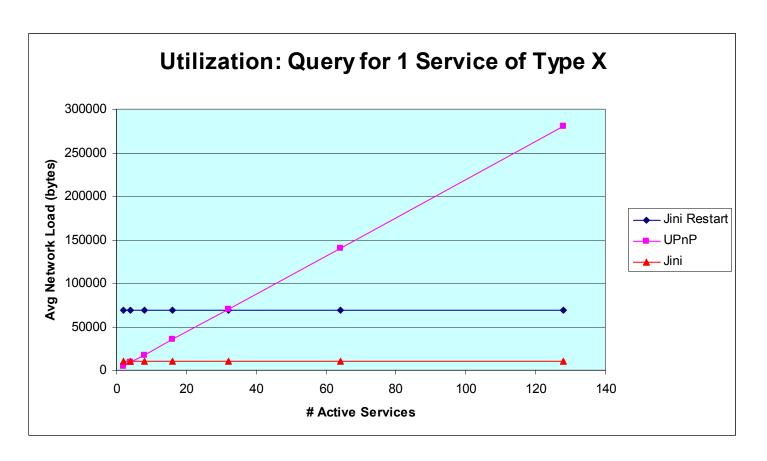
Performance of Query for Specific Service: "Find service X"



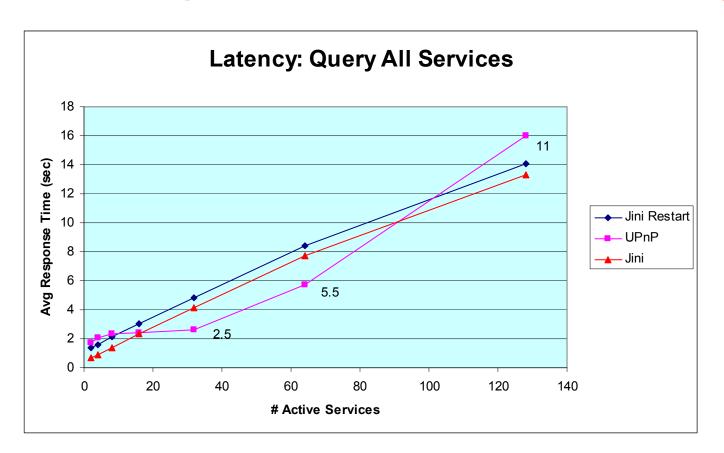
Performance of Query for Specific Service: "Find service X"



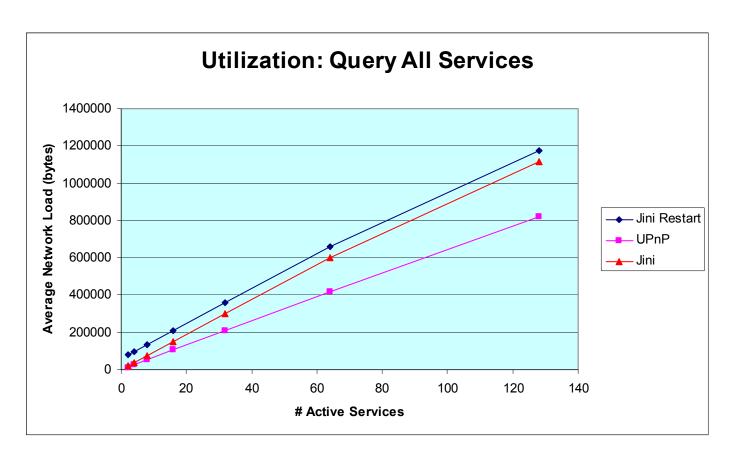
Performance of "anycast" Query: "Find one instance of type X"



Performance of "anycast" Query: "Find one instance of type X"



Performance of Service/Device Poll: "Find all active services"



Performance of Service/Device Poll: "Find all active services"